

Software System Design and Implementation

Admin & Motivation & Some History

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Admin

- **Course website** is the main source of information for this course:
 - www.cse.unsw.edu.au/~cs3141
- **Piazza as course forum**
 - you should receive invitation shortly
- **Consultation time**
 - Monday after the lecture - please send an email
 - otherwise, by appointment

Lectures

- Monday lectures will start at **9:15** and we skip the break
- Lectures will be recorded via Echo
- We're using 'OneNote' to share class whiteboard content
 - link and invitation coming soon

Assessment

- No tutes, but regular exercises and quizzes
 - to practice and deepen understanding of lecture content
- 2 Assignments
- Final exam
 - harmonic mean between class mark and exam
 - need to pass exam (40%) to pass the course

Motivation

What this course is all about?

Software affects our lives everywhere

- Mobile phones
- Social networks
- Cars & public transport
- Multimedia devices
- Medical systems
- Crypto currency
- Obviously, on computers at home and at work



Photo by Prosper973 - <http://flic.kr/p/6pECyk>



Photo by brykmantra - <http://flic.kr/p/7MrFA>



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Software is often mission critical.

CAN CLASS B

- 1 SAM/SRB Fahrer
- 2 SAM/SRB Beifahrer
- 3 SAM/SRB Heck 1
- 4 SAM/SRB Heck 2
- 5 Sitzsteuergerät Fahrer
- 6 Sitzsteuergerät Beifahrer
- 7 Sitzsteuergerät hinten links
- 8 Sitzsteuergerät hinten rechts
- 9 Türsteuergerät vorne Fahrerseite
- 10 Türsteuergerät vorne Beifahrerseite
- 11 Türsteuergerät hinten Fahrerseite
- 12 Türsteuergerät hinten Beifahrerseite
- 13 Steuergerät Trennwand
- 14 Dachbedieneinheit
- 15 Dachknoten Mitte (DKM)
- 16 Vorderes-Becken-Feld (VBF)
- 17 Hinteres-Becken-Feld (HBF)
- 18 Elektronisches Zündschloss (EZS)
- 19 Kombiinstrument
- 20 Mantelohrmodul
- 21 Frontklimatisierung
- 22 Fondklimatisierung
- 24 Audiogateway

CAN CLASS C

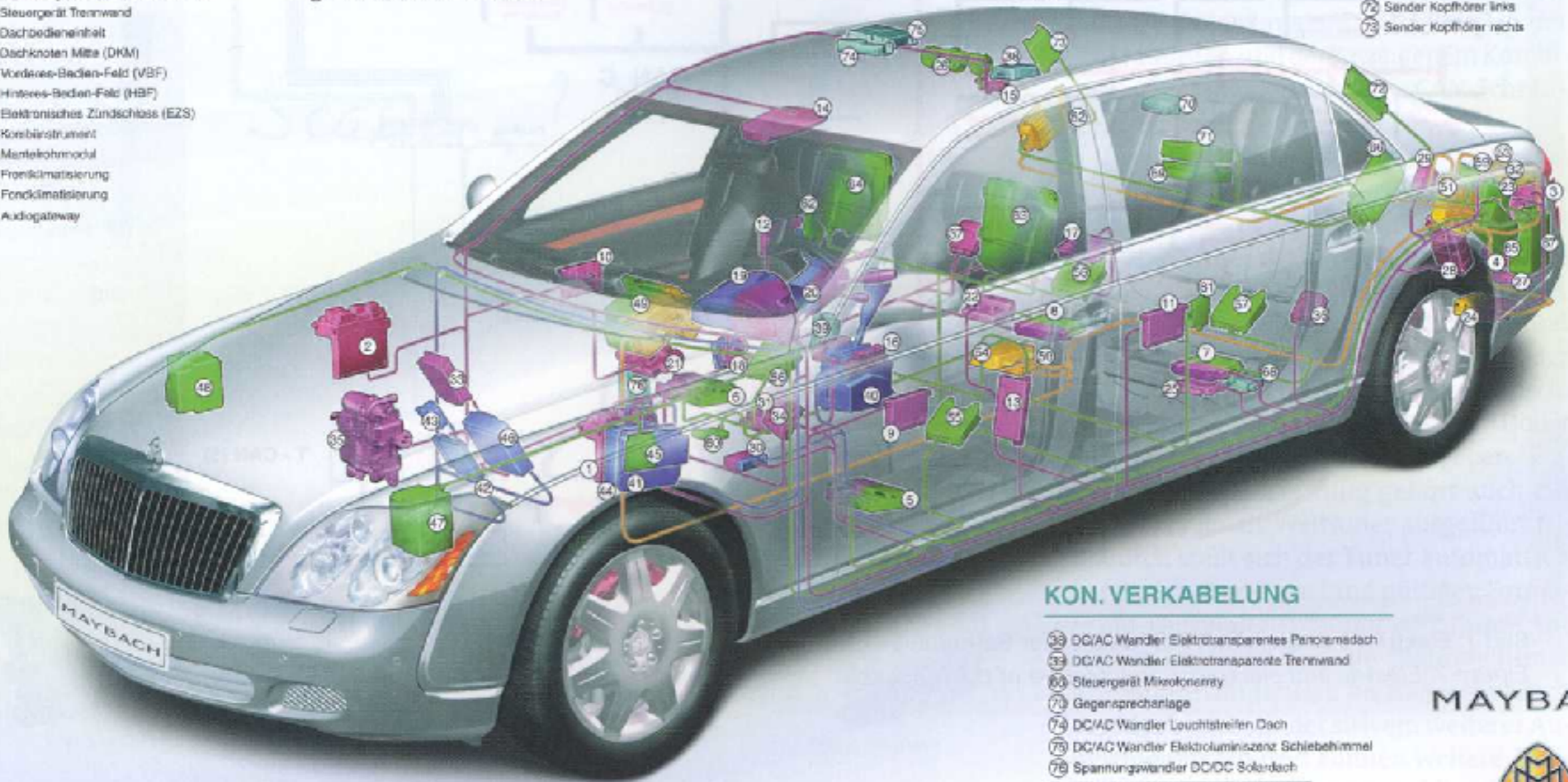
- 25 Parktronicasystem (PTS)
- 27 Reifendruckkontrolle (RDK)
- 28 Pneumatische Steuerinheit (PSE)
- 29 Heckdeckelfernschlüsselöffnung
- 30 Zentrales Gateway
- 31 Airbag-SG (Armada)
- 32 Multifunktionssteuergerät (MSS)
- 33 Bordnetz Steuergerät
- 34 Wandler Lenkradheizung
- 35 Standheizung
- 36 Türzuziehung hinten Fahrerseite
- 37 Türzuziehung hinten Beifahrerseite
- 38 Elektronisches Zündschloss (EZS)
- 39 Kombiinstrument
- 40 Mantelohrmodul
- 41 Zentrales Gateway
- 42 Elektronisches Wählhebelmodul
- 43 Luftfederung (SLF)
- 44 DISTRONIC (DTR)
- 45 Leuchtelektroregulierung
- 46 Motorelektronik (ME)
- 47 Sensotronic Brake System (FSG)
- 48 Elektronische Getriebe-Steuerung

MOST-BUS

- 24 Audiogateway
- 48 Headunit
- 50 Steuergerät Sprachbedienung
- 51 TV-Tuner MOST
- 52 Soundverstärker
- 53 Navigationsrechner
- 54 Kommunikationsplattform (CP1)

PRIVATE-BUS

- 5 Sitzsteuergerät Fahrer
- 6 Sitzsteuergerät Beifahrer
- 7 Sitzsteuergerät hinten links
- 8 Sitzsteuergerät hinten rechts
- 23 TV-Tuner CAN
- 26 Dachinstrument
- 45 Sensotronic Brake System (FSG)
- 47 Sensotronic Brake System (ASG1)
- 48 Sensotronic Brake System (ASG 2)
- 55 Multikonturlehne vorne links
- 56 Multikonturlehne vorne rechts
- 57 Multikonturlehne hinten links
- 58 Multikonturlehne hinten rechts
- 59 Keyless Go Heckmodul
- 60 Keyless Go Innenraummodul
- 61 Keyless Go Tür hinten links
- 62 Keyless Go Tür hinten rechts
- 63 Fondbildschirm links
- 64 Fondbildschirm rechts
- 65 Kommunikationsplattform Fond (CP2)
- 66 Surround Amplifier
- 67 Audio Video Controller
- 68 CD-Wechsler
- 71 DVD Spieler
- 72 Sender Kopfhörer links
- 73 Sender Kopfhörer rechts



KON. VERKABELUNG

- 33 DC/AC Wandler Elektrotransparentes Panoramadach
- 34 DC/AC Wandler Elektrotransparente Trennwand
- 60 Steuergerät Mikroklima
- 71 Gegensprechanlage
- 74 DC/AC Wandler Lichtstreifen Dach
- 75 DC/AC Wandler Elektrolumineszenz Schliebehimmel
- 76 Spannungswandler DC/DC Solardach

Σ aller Steuergeräte: 76

MAYBACH



Modern cars are computer networks on wheels

So, why is software still so unreliable?

Software must be of high quality:
correct, safe & secure

Software must be developed
with low effort: cheap & quickly

products include increasing amounts of software
shouldn't get more expensive due to
software

releases shouldn't be delayed due to software



Software defects affect many
activities

increasingly used in safety and
security critical applications

Financial success often depends on software
success

Sometimes we can sacrifice one for the other

- Computer games: **development effort is key**

- ▶ Correctness and safety are secondary
- ▶ Nobody notices a few pixel with the wrong colour
- ▶ or if the game physics isn't quite accurate....



Sometimes we can sacrifice one for the other

- Flight management system
(correctness is key)

- ▶ Safety and security is an overriding concern
- ▶ Defects are very costly and may harm human life

New York Times:

Airline Blames Bad Software in San Francisco Crash

By MATTHEW L. WALD MARCH 31, 2014



Usually, we need a balance

- Consider testing for a mass market product
 - ▶ If we test for a very long time with many testers
 - the product will be expensive and we ship too late, but with few defects
 - ▶ If we do very little testing
 - the product will be cheaper and ship early, but be riddled with bugs
- In practice, you want to ship when the remaining bugs are unlikely to affect users (sales?) very much

Implications

1. We need to be able to **trade quality for reduced effort**
 - ▶ To be broadly applicable, an approach to software design and implementation must support this trade off
2. We ideally want to **increase quality, while reducing effort**
 - ▶ We seek novel approaches that solve both problems at once

Produce better software with less effort

- Better software

- ▶ Software that has fewer defects (including security defects)
- ▶ Software that is more usable — we won't talk about usability in this course

- Less effort

- ▶ Shorter development time
- ▶ Fewer programmers
- ▶ Less-specialised programmers

The core theme of this course

How can mathematical tools & concepts help to produce better software with less effort?

Why are mathematical tools interesting?

- Mature engineering disciplines are based on mathematical foundations
 - ▶ We rarely *guess* the input current to a semiconductor
 - ▶ Would you build a bridge and then *test* whether it can hold its load?
- They enable a qualitative change in software development
 - ▶ We want to increase quality, while reducing effort
 - ▶ High-assurance software requires proof

What mathematical tools?

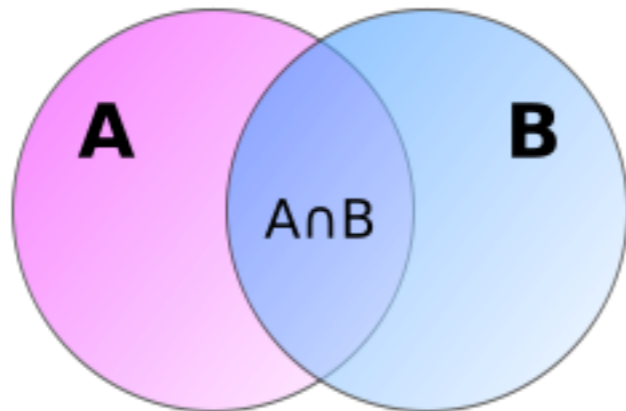
- Tools to reason about specifications and programs
 - ▶ Theorem provers & proof assistants
 - ▶ Static analysis tools & type checkers
- Tools to transform and refine programs
 - ▶ Meta programming & generative programming
 - ▶ Rewriting tools
- Advanced programming languages and environments

What is the basis of these tools?

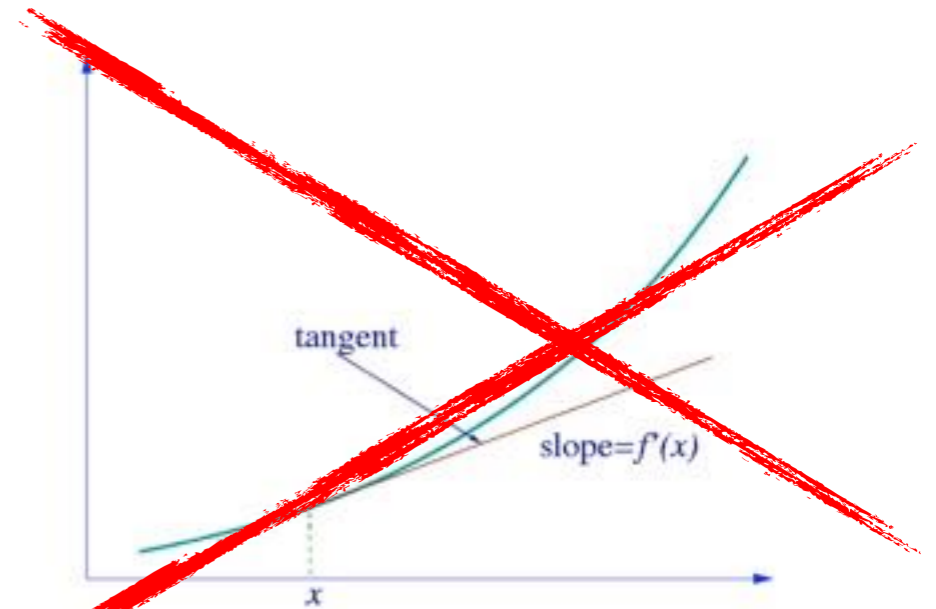
**The basis for reasoning
about programs**

$\forall x.P(x)$

Mathematical logic



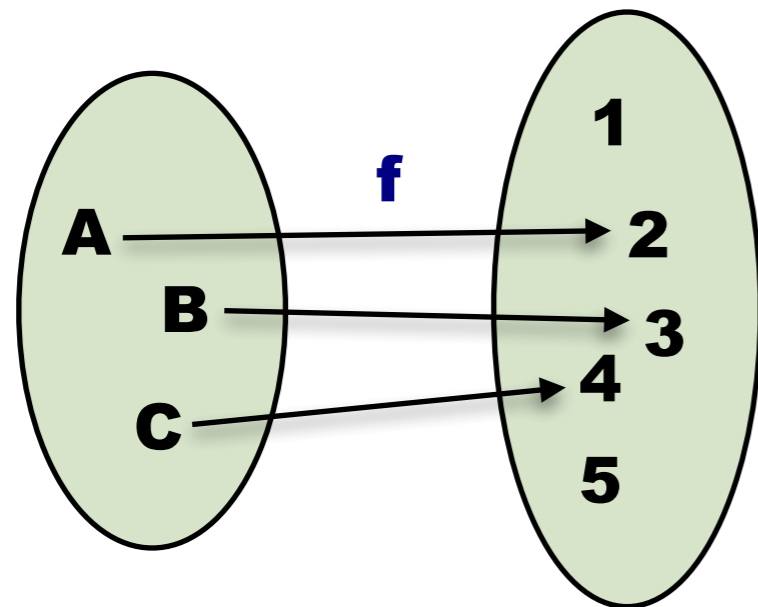
Discrete mathematics



Calculus

Discrete mathematics

- Study of structures that are fundamentally discrete rather than continuous
- In a discrete space, individual points are isolated from each other — consider, natural numbers versus real numbers



functions in computing

state space of a computer

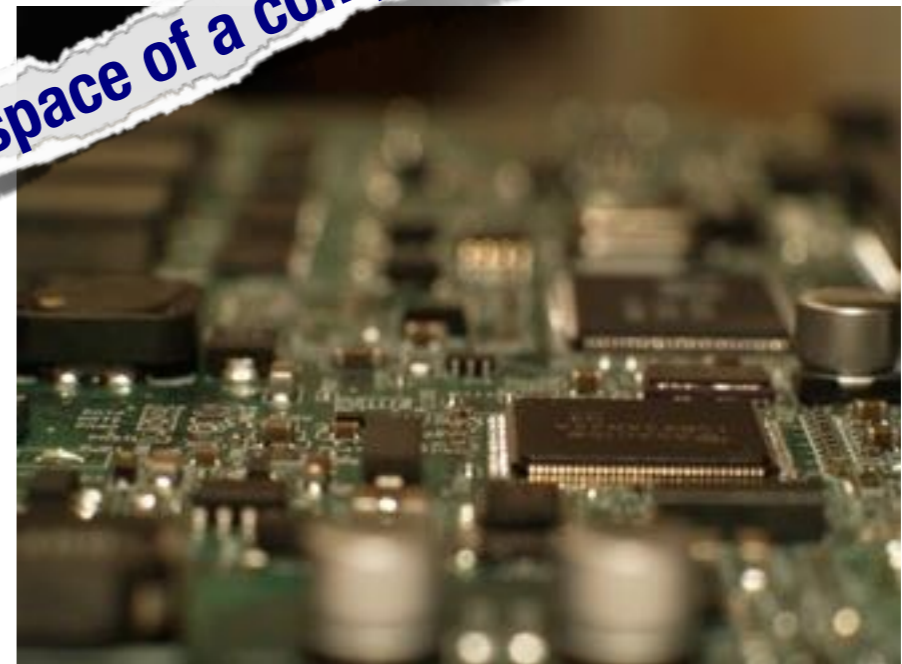


Photo by Un ragazzo chiamato Bi - <http://flic.kr/p/4NFrtx>

Mathematical logic

- Foundation of reasoning about computation, languages, and programs
 - ▶ Mathematical basis for software design and implementation
- Originally arose out of the study of the foundations of mathematics
 - ▶ On what ultimate basis can mathematical statements be called true?
 - ▶ The foundations of mathematics were heavily disputed in the early 20th century - mainly intuitionists vs formalist

A little bit of history

- 1900: Hilbert's Problems

- ▶ (Initially) ten open problems of mathematics
- ▶ 2nd problem:

Prove that the axioms of arithmetic are consistent



- 1920s: Hilbert's program

- ▶ Base all of mathematics on a finite, consistent set of axioms
- ▶ Show that they are complete, consistent

$$0 \in \mathbb{N}$$

$$x \in \mathbb{N} \Rightarrow s(x) \in \mathbb{N}$$

$$x \in \mathbb{N} \Rightarrow s(x) \neq 0$$

$$s(x) = s(y) \Rightarrow x = y$$

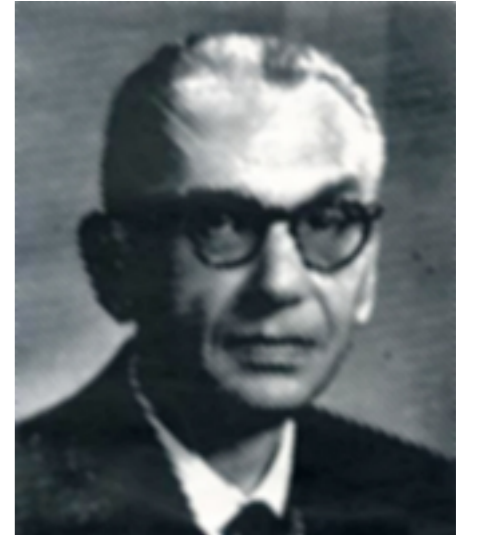
$$0 \in \mathbb{X}, (x \in \mathbb{N} \Rightarrow (x \in \mathbb{X} \Rightarrow s(x) \in \mathbb{X})) \Rightarrow \mathbb{N} \subseteq \mathbb{X}$$

- 1928: Hilbert's Entscheidungsproblem (decision problem)

- ▶ given a statement in first order logic, is the statement provable?

A little bit of history

- 1931: Gödel's incompleteness theorems
 - ▶ Consistent, computable axiom set, covering arithmetic, can never be complete
 - ▶ Such a system can't even prove it's own consistency



- Adapted form of Hilbert's program
 - ▶ Instead of formalising all of mathematics, formalise the important parts
- There are many consistent and complete logics
 - ▶ Gödel showed that **first-order logic** is complete

- 1928: Hilbert's Entscheidungsproblem ("decision problem")
 - ▶ Is there an algorithm that, given a mathematical statement (in a formal language), will tell us whether the statement is true or false?

- Church (1936) & Turing's (1937) Church-Turing Theorem

- ▶ Showed that such an algorithm is impossible



- Church-Turing Thesis

- ▶ Recursive functions, Turing machines & lambda calculus are equivalent

Halting Problem

- **Idea:** assume we have a total function

`stops (p, i):` returns True if program `p` applied to input `i` stops, false otherwise

`f (i) = if stops (i ,i)
 then loop forever
 else 1`

`f (f) ?`

Lambda calculus

- Minimal calculus that can express all computable functions
 - ▶ Only variables, function abstraction, and function application
- The basis for many theorem provers
- The foundation of all functional programming languages
 - ▶ Extend the lambda calculus with explicit data structures
 - ▶ Add syntactic sugar to make it more convenient

Background

- Very simple, but Turing-complete (Church-Turing thesis)
 - ▶ Pure lambda calculus has three constructs and two rewrite rules
 - ▶ It's **compositional** (i.e., it's highly modular)
- Today there are many different flavours
 - ▶ With and without types, with many extensions, and so on
- Lambda abstractions in C#, C++1x, Objective-C (as blocks), Swift & FP languages

Syntax

- The pure λ -calculus has three syntactic forms
 - ▶ **Variables:** x, y, z, \dots
 - ▶ **Lambda abstraction:** $\lambda x.M$
 - ▶ **Function application:** MN

Reduction rules

- Alpha conversion:

- ▶ We may change the name of any **bound** variable

- ▶ For example, we may convert $\lambda x.x$ to $\lambda y.y$

- ▶ A bound variable is one named in an enclosing lambda abstraction

- ▶ For example, here x is bound, but y is **free**: $\lambda x.xy$

Reduction rules

- Beta reduction:

- ▶ We can **reduce** a term of the form $(\lambda x.M)N$ to $M[N/x]$

- ▶ Where $M[N/x]$ means to replace any x in M by N

- ▶ The latter process is called **substitution**

- ▶ For example, we have

$$(\lambda x.x)y \longrightarrow y$$

$$(\lambda x.\lambda y.x)(\lambda z.z)vw \longrightarrow (\lambda y.\lambda z.z)vw \longrightarrow (\lambda z.z)w \longrightarrow w$$

execution of a lambda term

Church Encoding

- How can simple data types and operations on them be encoded in the lambda calculus
 - boolean values?
 - natural numbers?

Course contents

- Logical program properties
 - ▶ Basic logic & formal properties
 - ▶ Property-based testing
- Types help to design programs
 - ▶ Program properties as types
 - ▶ Types guide the design
 - ▶ Encapsulating properties
- Types help to implement programs
 - ▶ Types imply programs
 - ▶ Types control effects
 - ▶ Types prevent defects
- Effect control helps with parallelism
 - ▶ Effects interfere with concurrency & parallelism

Haskell



- A practical, strongly-typed functional programming language
 - ▶ Widely used in research, industry & education
 - ▶ Mature, highly optimising compiler with interactive environment
 - ▶ Over thousands of open-source libraries and tools
- Named after the logician Haskell B. Curry

<http://haskell.org/>

Why Haskell?

- Functional languages are based on the lambda calculus
 - ▶ Semantics of programs is fairly precisely defined
 - ▶ This simplifies formal reasoning about these programs
- Functional languages can dramatically increase productivity
 - ▶ Factor of four has been cited for Erlang versus C++
- Haskell has a very sophisticated type system
- Haskell has controlled effects